

# Tree Automata and Applications

## Exercise session 1

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### Exercise 1 - First constructions of Tree Automata

Let  $\mathcal{F} = \{f(2), g(1), a(0)\}$ . Give a DFTA and a top-down DFTA for the set  $G(t)$  of ground instances of the term  $t = f(f(a, x), g(y))$  which is defined by:

$$G(t) = \{f(f(a, u), g(v)) \mid u, v \in T(\mathcal{F})\}$$

### Exercise 2 - What is recognizable by a tree automaton?

Are the following tree languages recognizable by a bottom-up tree automaton?

1.  $\mathcal{F} = \{g(1), a(0)\}$  and  $L$  the set of ground terms of even height.
2.  $\mathcal{F} = \{f(2), g(1), a(0)\}$  and  $L$  the set of ground terms of even height.

### Exercise 3 - Bottom-up vs Top-down

1. Recall why bottom-up NFTAs, bottom-up DFTAs and top-down NFTAs have the same expressiveness.
2. Let  $\mathcal{F} = \{f(2), g(1), a(0)\}$ . Give a DFTA and a top-down NFTA for the set  $M(t)$  of terms which have a ground instance of the term  $t = f(a, g(x))$  as a subterm, i.e.

$$\{C[f(a, g(u))] \mid C \in \mathcal{C}(\mathcal{F}), u \in T(\mathcal{F})\}.$$

3. Show that NFTAs and top-down DFTAs do not have the same expressiveness.

### Homework: Satisfiability

Let  $\mathcal{F} = \{\text{and}(2), \text{or}(2), \text{not}(1), 0(0), 1(0), x(0)\}$ . A ground term over  $\mathcal{F}$  can be viewed as a boolean formula over  $x$ .

1. Give an NFTA which recognizes the set of satisfiable boolean formulae over  $x$ . Show that it indeed does.

Let  $\mathcal{F} = \{\text{and}(2), \text{or}(2), \text{not}(1), 0(0), 1(0), x_1(0), \dots, x_{n(0)}\}$ , i.e. we now handle  $n$  variables instead of a single one. The same variable may appear several times in a formula, and should be evaluated consistently.

2. Give an NFTA which recognizes the set of satisfiable boolean formulae over  $x_1, \dots, x_n$ . Show that it does.